

2-STRONG PRODUCT AND CONNECTEDNESS

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Abstract: Among all the product of two graphs G and H , the Cartesian product $G \times H$, tensor product $G \otimes H$ and strong product $G \boxtimes H$ are very well known and studied in detail. Recently, Cartesian product and tensor product were generalized by defining 2-Cartesian product and 2-tensor product and their properties were studied. In this paper, we have generalized strong product of two graphs G and H by defining 2-strong product $G \boxtimes_2 H$ and studied some basic graph parameters like connectedness and distance.

Keywords and Phrases: Strong product of graphs, connected graph, distance, diameter, eccentricity, radius.

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1. Introduction and Preliminaries

Product of two graphs G and H has been defined in many different ways in literature. Among all the products, Cartesian product $G \times H$, tensor product

$G \otimes H$, strong product $G \boxtimes H$ and lexicographic product $G \circ H$ are studied in detail ([5] & [6]). Cartesian product and tensor product have been generalized by defining 2-Cartesian product $G \times_2 H$ and 2-tensor product $G \otimes_2 H$ in [2] & [3] and basic parameters of both the graphs have been obtained in terms of the parameters of G and H .

In this paper, we generalize strong product of two graphs G and H by defining 2-strong product $G \boxtimes_2 H$ and study some basic graph parameters.

For any graph to be called connected, every vertex pair in it should be connected by a path. Distance between two vertices x and x' in graph G is defined as the length of shortest path between x and x' . For rest of the basic definitions in graph, we refer [4].

Definition 1.1. Let $G = (V(G), E(G))$ and $H = (V(H), E(H))$ be two connected graphs. Then 2-strong product of G and H is a graph denoted by $G \boxtimes_2 H$ with vertex set as $V(G \boxtimes_2 H) = V(G) \times V(H)$ and edge set as $E(G \boxtimes_2 H) = \{(x, y)(x', y') : (x = x' \text{ \& } d_H(y, y') = 2) \text{ or } (d_G(x, x') = 2 \text{ \& } y = y') \text{ or } (d_G(x, x') = 2 \text{ \& } d_H(y, y') = 2)\}$. Equivalently, $d_G(x, x') \in \{0, 2\}$ and $d_H(y, y') \in \{0, 2\}$.

If we replace 2 by 1, we get the definition of usual strong product $G \boxtimes H$.

Remark 1.2.

$$1. |E(G \boxtimes_2 H)| = |V(G)||N^2(H)| + |V(H)||N^2(G)| + 2|N^2(G)||N^2(H)|,$$

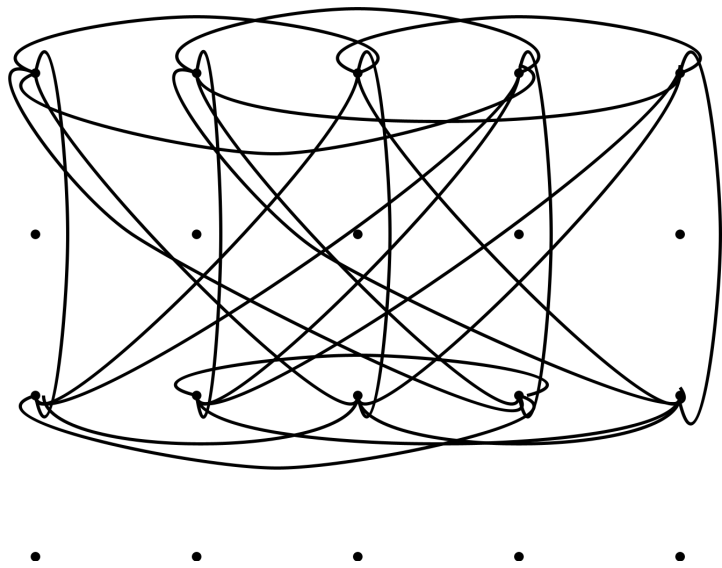
where $N^2(G)$ denotes the collection of all distinct pair of vertices which are at distance 2. Hence, we always assume that $N^2(G)$ and $N^2(H)$ are non-empty sets. Otherwise, we get $G \boxtimes_2 H$ as null graph.

$$2. G \boxtimes_2 H \cong H \boxtimes_2 G.$$

$$3. G \boxtimes_2 H = (G \times_2 H) \cup (G \otimes_2 H), \text{ where } G \times_2 H \text{ is the 2-Cartesian product of } G \text{ and } H, G \otimes_2 H \text{ is the 2-tensor product of } G \text{ and } H \text{ and } \cup \text{ denotes the edge disjoint union of } G \times_2 H \text{ and } G \otimes_2 H \text{ ([2], [3])}.$$

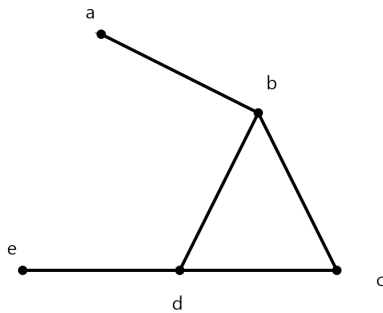
Example 1.3.

1. $P_4 \boxtimes_2 C_4$ has four components as K_4 .
2. One component of 2-strong product of two graphs P_4 and C_5 , $P_4 \boxtimes_2 C_5$ is shown below.

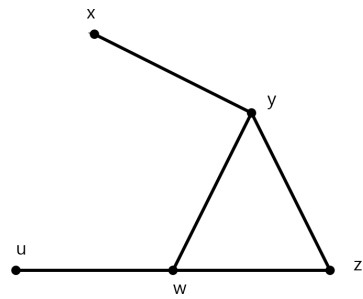


3. Let G & H be the graphs shown below.

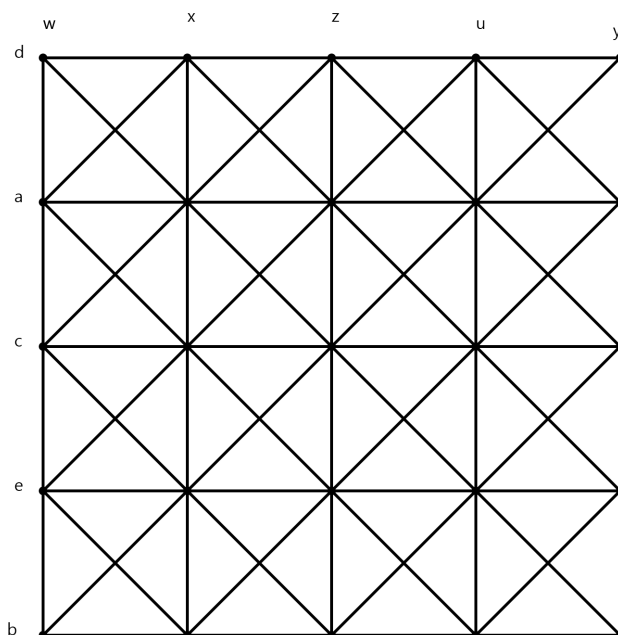
Figure 1: G



H



Then, $G \boxtimes_2 H$ is shown below.

Figure 2: $G \boxtimes_2 H$ 

First, we obtain basic graph parameters of $G \boxtimes_2 H$ such as degree of vertex, regularity and Eulerian property.

Proposition 1.4. *Let $G = (V(G), E(G))$ and $H = (V(H), E(H))$ be two graphs. Then, for $(x, y) \in V(G \boxtimes_2 H)$,*

$$\deg_{G \boxtimes_2 H}((x, y)) = \deg_2(x) + \deg_2(y) + \deg_2(x) \deg_2(y),$$

where $\deg_2(x) = |N^2(x)|$ and $N^2(x) = \{x' \in V(G) : d_G(x, x') = 2\}$.

Proof. Let $(x, y) \in V(G \boxtimes_2 H)$. Then all vertices (x, y') , where $y' \in V(H)$ with $d_H(y, y') = 2$ are adjacent to (x, y) . Note that there are $\deg_2(y)$ such vertices. Similarly, all vertices (x', y) , where $x' \in V(G)$ with $d_G(x, x') = 2$ are adjacent to (x, y) . There are $\deg_2(x)$ such vertices. Also, note that $(x', y') \in V(G \boxtimes_2 H)$ where $d_G(x, x') = 2$ and $d_H(y, y') = 2$ is adjacent to (x, y) and there are $\deg_2(x)\deg_2(y)$ such vertices. Thus, degree of (x, y) in $G \boxtimes_2 H$ is $\deg_2(x) + \deg_2(y) + \deg_2(x)\deg_2(y)$.

Now, we obtain a sufficient condition of regularity in $G \boxtimes_2 H$, using the following definition of second regular graph given in [7].

Definition [7]. A graph $G = (V(G), E(G))$ is said to be second regular with

regularity k if $\deg_2(x) = k$, for all $x \in V(G)$.

Corollary 1.5. *If graphs G and H are second regular graphs with regularity k_1 and k_2 respectively, then $G \boxtimes_2 H$ is $k_1 + k_2 + k_1k_2$ regular graph.*

In general, if G and H are Euler graphs, then $G \boxtimes_2 H$ may not be an Euler graph.

Example 1.6. Let $G = H$ be the graph shown below.

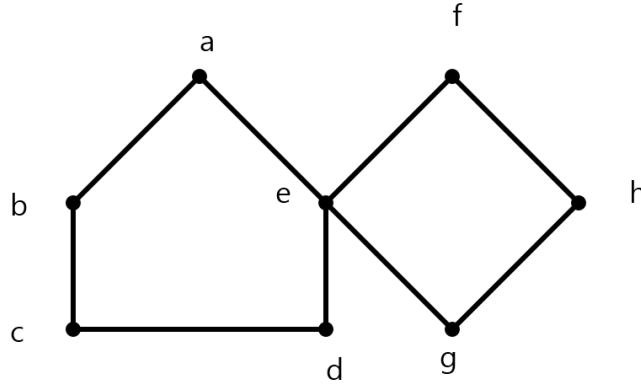


Figure 3: $G = H$

Then, G and H are Euler graphs, but $G \boxtimes_2 H$ is not Euler, as $\deg_{G \boxtimes_2 H}(e, e)$ is odd, using Proposition 1.4.

Now, we give a characterization for $G \boxtimes_2 H$ to be Euler graph, in terms of 2 degree of a vertex.

Theorem 1.7. *Let $G = (V(G), E(G))$ and $H = (V(H), E(H))$ be two connected graphs, such that $G \boxtimes_2 H$ is connected. Then, $\deg_2(x)$ is even for every $x \in V(G)$ and $\deg_2(y)$ is also even for every $y \in V(H)$ if and only if $G \boxtimes_2 H$ is an Euler graph.*

Proof. Follows from Proposition 1.4.

Remark 1.8. *Let $G = (V(G), E(G))$ and $H = (V(H), E(H))$ be two graphs with $N^2(x) \neq \emptyset$ for some $x \in V(G)$ and $N^2(y) \neq \emptyset$ for some $y \in V(H)$. Then K_4 is a subgraph of $G \boxtimes_2 H$ and hence $G \boxtimes_2 H$ is a non-bipartite graph.*

2. Distance and Connectedness

In this section, we obtain the distance formula and discuss connectedness of $G \boxtimes_2 H$.

In usual strong product $G \boxtimes H$ of two graphs G and H , distance formula is

given as follows:

Theorem. [5] For two vertices (x, y) and (x', y') in $G \boxtimes H$,

$$d_{G \boxtimes H}((x, y), (x', y')) = \max\{d_G(x, x'), d_H(y, y')\}.$$

To obtain distance formula for $G \boxtimes_2 H$, we shall need the following definition.

Definition. [7] Let G be a connected graph with $N^2(u) \neq \emptyset$, $\forall u \in V(G)$ and let $x, x' \in V(G)$. Then, $d'_G(x, x')$ is defined as the length of the shortest walk $W : x = w_0 \rightarrow w_1 \rightarrow \dots \rightarrow w_{2k} = x'$ between x and x' of the form $2k$ ($k \in \mathbb{N}$) in which $d(w_i, w_{i+2}) = 2$, $i = 0, 2, 4, \dots, 2k - 2$.

If there is no such walk between x and x' , then $d'_G(x, x') = \infty$.

Remark 2.1.

1. For any $x, x' \in V(G)$, $d_G(x, x') \leq d'_G(x, x')$.
2. For $x, x' \in V(G)$, if $d_G(x, x')$ is even, then $d_G(x, x') = d'_G(x, x')$ and if $d_G(x, x')$ is odd, then $d_G(x, x') < d'_G(x, x')$.
3. For a connected bipartite graph G , if x and x' are in same partite sets of G , then $d_G(x, x')$ is even and hence $d'_G(x, x') = d_G(x, x') < \infty$, whereas, if x and x' are in different partite sets, then $d_G(x, x') < \infty$ but $d'_G(x, x') = \infty$.

Example 2.2. For $G = C_{2n+1}$, if $d_G(x, x')$ is odd, then $d'_G(x, x') = (2n + 1) - d_G(x, x') > d_G(x, x')$.

Proposition 2.3. For a connected, non bipartite and triangle free graph G , $d'_G(x, x') < \infty$ for every $x, x' \in V(G)$.

Proof. Let $x, x' \in V(G)$ and C be an odd cycle in G . If $d_G(x, x')$ is even, then $d'_G(x, x') = d_G(x, x') < \infty$. Now, if $d_G(x, x')$ is odd, then traversing the cycle C , there is an even length walk, say W between x and x' . Let $W : x = u_0 \rightarrow u_1 \rightarrow u_2 \rightarrow \dots \rightarrow u_{2k} = x'$ be the walk. Then, note that, if $d_G(u_i, u_{i+2}) = 1$ for some even $i \leq 2k$, then u_i, u_{i+1} & u_{i+2} forms a triangle in G , which is not possible. Hence, $d_G(u_i, u_{i+2}) = 2$, $i = 0, 2, 4, \dots, 2k - 2$ and therefore, $d'_G(x, x') \leq 2k < \infty$.

Using this definition of $d'(x, x')$, we obtained the distance formula for two vertices in $G \boxtimes_2 H$.

Theorem 2.4. Let G and H be two graphs with $N^2(x) \neq \emptyset$ for all $x \in V(G) \cup V(H)$. Then for $(x, y), (x', y') \in V(G \boxtimes_2 H)$,

$$d_{G \boxtimes_2 H}((x, y), (x', y')) = \max\left\{\frac{d'_G(x, x')}{2}, \frac{d'_H(y, y')}{2}\right\}.$$

Proof. Suppose $d'_G(x, x') = 2m < \infty$ and $d'_H(y, y') = 2n < \infty$. Consider, the following two walks corresponding to $d'_G(x, x')$ and $d'_H(y, y')$ with $d_G(x_i, x_{i+2}) = 2 = d_H(y_i, y_{i+2})$ ($i = 0, 2, 4, \dots$):

$$\begin{aligned} P : x = x_0 \rightarrow x_1 \rightarrow x_2 \rightarrow \dots \rightarrow x_{2m-1} \rightarrow x_{2m} = x' \text{ \& } \\ Q : y = y_0 \rightarrow y_1 \rightarrow y_2 \rightarrow \dots \rightarrow y_{2n-1} \rightarrow y_{2n} = y' \end{aligned}$$

Without loss of generality, assume that $m \geq n$. Now,

$(x, y) \rightarrow (x_2, y_2) \rightarrow (x_4, y_4) \rightarrow \dots \rightarrow (x_{2n}, y_{2n}) \rightarrow (x_{2n+2}, y_{2n}) \rightarrow \dots \rightarrow (x', y')$ is a walk in $G \boxtimes_2 H$ of length m as $d_G(x_i, x_{i+2}) = d_H(y_i, y_{i+2}) = 2$. Thus,

$$d_{G \boxtimes_2 H}((x, y), (x', y')) \leq m = \frac{d'_G(x, x')}{2} = \max\left\{\frac{d'_G(x, x')}{2}, \frac{d'_H(y, y')}{2}\right\} < \infty.$$

Conversely, let $d_{G \boxtimes_2 H}((x, y), (x', y')) = m < \infty$ and let

$P : (x, y) \rightarrow (x_1, y_1) \rightarrow (x_2, y_2) \rightarrow \dots \rightarrow (x_{m-1}, y_{m-1}) \rightarrow (x', y')$ be corresponding path in $G \boxtimes_2 H$. Now, consider the paths

$$\begin{aligned} P_G(P) : x = x_0 \rightarrow x_1 \rightarrow x_2 \rightarrow \dots x_{m-1} \rightarrow x_m = x' \text{ \& } \\ P_H(P) : y = y_0 \rightarrow y_1 \rightarrow y_2 \rightarrow \dots \rightarrow y_{m-1} \rightarrow y_m = y' \end{aligned}$$

with $d_G(x_i, x_{i+1}) = 0$ or 2 for every i and $d_H(y_i, y_{i+1}) = 0$ or 2 for every i .

Case 1. $x = x_1 = x_2 = \dots = x_{m-1} = x_m = x'$

In this case, $d_G(x, x') = 0$ and hence $d'_G(x, x') = 0$ by definition. Also, $y_i \neq y_j$, for any $1 \leq i, j \leq m$, ($i \neq j$) and hence $d_H(y_{i+1}, y_i) = 2$ ($i = 0, 2, \dots$). Now, as $d_H(y_{i+1}, y_i) = 2$ for every i , we have $y_i \rightarrow a_{i+1} \rightarrow y_{i+1}$ for every i , for some $a_{i+1} \in V(H)$. Thus, $y = y_0 \rightarrow a_1 \rightarrow y_1 \rightarrow a_2 \rightarrow y_2 \rightarrow \dots \rightarrow \dots a_k \rightarrow y_m = y'$ with $d_H(y_i, y_{i+1}) = 2$ for every i . Hence, we get $d'_H(y, y') \leq 2m$. Therefore, $\frac{d'_H(y, y')}{2} \leq m = d_{G \boxtimes_2 H}((x, y), (x', y'))$. Also, $\frac{d'_G(x, x')}{2} = 0$. Thus,

$$\max\left\{\frac{d'_H(y, y')}{2}, \frac{d'_G(x, x')}{2}\right\} \leq d_{G \boxtimes_2 H}((x, y), (x', y')).$$

Similarly, if $y = y_1 = y_2 = \dots y_{m-1} = y_m = y'$, then

$$\max\left\{\frac{d'_H(y, y')}{2}, \frac{d'_G(x, x')}{2}\right\} \leq d_{G \boxtimes_2 H}((x, y), (x', y')).$$

Case 2. There is atleast one i , such that $x_i \neq x$ and j such that $y_j \neq y$.

Note that in this case, from $P_G(P)$ and $P_H(P)$, we get two subsequences

$x = u_0 \rightarrow u_1 \rightarrow u_2 \rightarrow \dots u_{l-1} \rightarrow u_l = x'$ with $d_G(u_i, u_{i+1}) = 2$ for every i and

$y = v_0 \rightarrow v_1 \rightarrow v_2 \rightarrow \dots v_{k-1} \rightarrow v_k = y'$ with $d_H(v_i, v_{i+1}) = 2$ for every i .

Now as $d_G(u_i, u_{i+1}) = 2$ and $d_H(v_i, v_{i+1}) = 2$ for every i , there is $a_{i+1} \in V(G)$ and $b_{i+1} \in V(H)$ such that $u_i \rightarrow a_{i+1} \rightarrow u_{i+1}$ and $v_i \rightarrow b_{i+1} \rightarrow v_{i+1}$. Thus, we get two walks, $x = u_0 \rightarrow a_1 \rightarrow u_1 \rightarrow a_2 \rightarrow u_2 \rightarrow \dots a_l \rightarrow u_l = x'$ and $y = v_0 \rightarrow b_1 \rightarrow v_1 \rightarrow b_2 \rightarrow v_2 \rightarrow \dots \rightarrow b_k \rightarrow v_k = y'$. Thus $d'_G(x, x') \leq 2l$ and $d'_H(y, y') \leq 2k$.

Therefore, $\frac{d'_G(x, x')}{2} \leq l \leq m$ and $\frac{d'_H(y, y')}{2} \leq k \leq m$. Thus,
 $\max\{\frac{d'_G(x, x')}{2}, \frac{d'_H(y, y')}{2}\} \leq m$.

Now, suppose $d'_G(x, x') = \infty$ or $d'_H(y, y') = \infty$. If possible, suppose $d_{G \boxtimes_2 H}((x, y), (x', y')) = m < \infty$. Then, using arguments given in converse part, $d'_G(x, x') < \infty$ and $d'_H(y, y') < \infty$, which is a contradiction. Now, suppose $d_{G \boxtimes_2 H}((x, y), (x', y')) = \infty$. If possible, suppose, $d'_G(x, x') < \infty$ and $d'_H(y, y') < \infty$. Then, by arguments given in first part, $d_{G \boxtimes_2 H}((x, y), (x', y')) < \infty$. Thus, $d'_G(x, x') = \infty$ and $d'_H(y, y') = \infty$.

Now, we obtain the connectedness of $G \boxtimes_2 H$.

Theorem 2.5. *Let G and H be two connected graphs with $N^2(u) \neq \emptyset \forall u \in V(G) \cup V(H)$.*

1. *If G and H are both bipartite graphs, then $G \boxtimes_2 H$ has four components.*
2. *If one of the graph G or H is non-bipartite triangle free and other is bipartite, then $G \boxtimes_2 H$ has 2 components.*
3. *If G and H both are non-bipartite triangle free graphs, then $G \boxtimes_2 H$ is connected.*

Proof.

1. Let U_1, U_2 be two partite sets of G and V_1, V_2 be two partite sets of H and let $W_{ij} = U_i \times V_j$ for $1 \leq i, j \leq 2$. Now, if $(x, y), (x', y')$ are in different W_{ij} , then either x and x' are in different partite sets of G or y and y' are in different partite sets of H . Therefore, either $d'_G(x, x') = \infty$ or $d'_H(y, y') = \infty$ and hence, by above theorem, $d_{G \boxtimes_2 H}((x, y), (x', y')) = \infty$. Thus, $G \boxtimes_2 H$ has at least four components W_{ij} .

Now, if (x, y) and (x', y') are in same W_{ij} , ($i, j \in \{1, 2\}$), then x & x' are in same partite set of G and y & y' are in same partite set of H . Thus, $d'_G(x, x') < \infty$ and $d'_H(y, y') < \infty$ and hence, $d_{G \boxtimes_2 H}((x, y), (x', y')) < \infty$, i.e., (x, y) and (x', y') are connected by a path in $G \boxtimes_2 H$.

Therefore, there are exactly four components in $G \boxtimes_2 H$.

2. Let G be a bipartite graph and let U_1 and U_2 be two partite sets of G and let $W_i = U_i \times V(H)$ for $1 \leq i \leq 2$. If (x, y) and (x', y') are in different W_i , then x and x' are in different partite sets of G and hence, $d'_G(x, x') = \infty$. Therefore, $d_{G \boxtimes_2 H}((x, y), (x', y')) = \infty$ and so, there are at the least two components in $G \boxtimes_2 H$.

Further, if (x, y) and (x', y') are in same W_i , $1 \leq i \leq 2$, then x and x' are in same partite set and hence, $d'_G(x, x') < \infty$. Also, as H is non bipartite and triangle free, by Proposition 2.3, $d'_H(y, y') < \infty$. Thus, $d_{G \boxtimes_2 H}((x, y), (x', y')) < \infty$ and hence (x, y) and (x', y') are connected by a path. Therefore, there are two components in $G \boxtimes_2 H$.

3. Let $(x, y), (x', y') \in V(G \boxtimes_2 H)$. Then, by Proposition 2.3, $d'_G(x, x') < \infty$ and $d'_H(y, y') < \infty$. Thus, $d_{G \boxtimes_2 H}((x, y), (x', y')) < \infty$ and hence, (x, y) and (x', y') are connected by a path. Thus, $G \boxtimes_2 H$ is connected.

3. Diameter, Eccentricity and Radius

In this section, we obtain some more basic graph parameters which are depending on distance between two vertices. For example, diameter, eccentricity and radius.

Definition 2.1. Let $G = (V(G), E(G))$ be a graph. Then, we define

1. 2-diameter of G by $\max\{d'_G(x, y) : x, y \in V(G)\}$ and we denote it by $\text{diam}_2(G)$.
2. 2-eccentricity of a vertex x in G as $\max\{d'_G(x, x') : x' \in V(G)\}$ and we denote it by $e'(x)$.
3. 2-radius of G as $\min\{e'(x) : x \in V(G)\}$ and we denote it by $\text{rad}_2(G)$.
4. 2-closed neighborhood of a vertex x in G as $\{x' \in V(G) : d'_G(x, x') \leq i\}$ and we denote it by $N'_i[x, G]$.

If we replace $d'_G(x, x')$ by $d_G(x, x')$ in above definitions, we get the definition of usual diameter, eccentricity, radius and closed neighborhood.

Remark 3.2. In general, $\text{diam}(G)$ (or $\text{rad}(G)$) and $\text{diam}_2(G)$ (or $\text{rad}_2(G)$) are different. Note that, $\text{diam}_2(G)$ (or $\text{rad}_2(G)$) is always even and if $\text{diam}(G)$ (or $\text{rad}(G)$) is even, still $\text{diam}_2(G)$ (or $\text{rad}_2(G)$) may be different from $\text{diam}(G)$ (or $\text{rad}(G)$).

We obtained the following relation between diameter of graph $G \boxtimes_2 H$ and 2-diameter of G and H .

Theorem 3.3. Let G and H be two graphs. Then,

$$\text{diam}(G \boxtimes_2 H) = \max\left\{\frac{\text{diam}_2(G)}{2}, \frac{\text{diam}_2(H)}{2}\right\}.$$

Proof. Let $\text{diam}(G \boxtimes_2 H) = d$. Then, for any $x, x' \in V(G)$ and $y, y' \in V(H)$, $d_{G \boxtimes_2 H}((x, y), (x', y')) \leq d$. Therefore, by Theorem 2.4, $\max\left\{\frac{d'_G(x, x')}{2}, \frac{d'_H(y, y')}{2}\right\} \leq d$.

Thus, $d'_G(x, x') \leq 2d$ & $d'_H(y, y') \leq 2d$. But, as x, x' and y, y' are arbitrary, $diam_2(G) \leq 2d$ and $diam_2(H) \leq 2d$. Therefore, $\max\{\frac{diam_2(G)}{2}, \frac{diam_2(H)}{2}\} \leq d$.

Conversely, let $diam_2(G) = d_1$ and $diam_2(H) = d_2$. Then, for any $(x, y), (x', y') \in V(G \boxtimes_2 H)$, $d'_G(x, x') \leq d_1$ and $d'_H(y, y') \leq d_2$. Therefore, $d_{G \boxtimes_2 H}((x, y), (x', y')) = \max\{\frac{d'_G(x, x')}{2}, \frac{d'_H(y, y')}{2}\} \leq \max\{\frac{d_1}{2}, \frac{d_2}{2}\}$. But, as (x, y) and (x', y') are arbitrary, $diam(G \boxtimes_2 H) \leq \max\{\frac{d_1}{2}, \frac{d_2}{2}\}$.

By similar arguments, the following can be proved.

Theorem 3.4. Let $G = (V(G), E(G))$ and $H = (V(H), E(H))$ be two graphs.

1. For $(x, y) \in V(G \boxtimes_2 H)$, $ecc(x, y) = \max\{\frac{e'(x)}{2}, \frac{e'(y)}{2}\}$.
2. $rad(G \boxtimes_2 H) = \max\{\frac{rad_2(G)}{2}, \frac{rad_2(H)}{2}\}$.
3. For $(x, y) \in V(G \boxtimes_2 H)$, $N_i[(x, y), G \boxtimes_2 H] = N'_{2i}[x, G] \times N'_{2i}[y, H]$, where $N_i[x, G] = \{x' \in V(G) : d(x, x') \leq i\}$.

We consider the following non-bipartite graphs:

Definitions [1]

1. A wheel graph W_n is a graph with $n + 1$ vertices that contains C_n and one other vertex which is adjacent to every vertex of C_n . The vertex which is adjacent to every vertex is called center vertex.
2. The Helm graph H_n is the graph obtained from W_n by adding a pendent edge to each vertex of C_n in W_n .
3. The closed Helm graph CH_n is the graph obtained from H_n by adding edges between pendent vertices.

Example 3.5.

Graph	$diam(G)$	$diam_2(G)$	$rad(G)$	$rad_2(G)$
C_{2n+1} ($n > 1$)	n	$2n$	n	$2n$
H_n ($n \geq 4$)	4	6	2	4
CH_n ($n \geq 4$)	4	6	2	4
P_n ($n > 2$)	$n - 1$	∞	$[n]$	∞
$K_{m,n}$ ($m, n > 2, m \neq n$)	2	∞	2	∞
W_n ($n > 4$)	2	∞	2	∞

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